



IN APPLIED SCIENCE & ENGINEERING TECHNOLOGY

Volume: 6 Issue: IV Month of publication: April 2018

DOI: http://doi.org/10.22214/ijraset.2018.4784

www.ijraset.com

Call: 🕥 08813907089 🔰 E-mail ID: ijraset@gmail.com



LICT Subdivision Double Domination in Graphs

M. H. Muddebihal¹, Suhas P. Gade²

¹Department of Mathematics Gulbarga University, Kulburgi-585106, Karnataka, India. ²Department of Mathematics, Sangameshwar College, Solapur – 413001, Maharashtra, India.

Abstract: Let S(G) be the subdivision graph of G. The lict graph of n[S(G)] of S(G) is a graph whose vertex set is the union of the set of edges and set of cutvertices of S(G) in which two vertices adjacent if and only if the corresponding members are adjacent or incident. A subset D^d of V[n(S(G))] is double dominating set of n[S(G)] if for every vertex $v \in V[n(S(G))]$, $|N[v] \cap$ $D^d| \ge 2$, that is v is in D^d and has at least one neighbour in D^d or v is in $V[n(S(G))] - D^d$ and has at least two neighbours in D^d . The lict subdivision double dominating number $\gamma_{ddns}(G)$ is a minimum cardinality of the lict subdivision double dominating set of G and is denoted by $\gamma_{ddns}(G)$. In this paper, we establish some sharp bounds for $\gamma_{ddns}(G)$. Also some upper and lower bounds on $\gamma_{ddns}(G)$ in terms of the vertices, edges and other different parameters of Gand not in terms of the element of n[S(G)]. Further, its relation with other different dominating parameters is also obtained. Subject classification number: AMS - 05C69, 05C70.

Keyword: Lict subdivision graph/ Dominating set/Double domination

INTRODUCTION

I.

In this paper, all the graphs considered here are simple, finite, non-trivial, undirected and connected. The vertex set and edge set of graph G are denoted by V(G) = p and E(G) = q respectively. Terms not defined here are used in the sense of Harary [1]. The neighbourhood of a vertex $v \in V$ is defined by $N(v) = \{u \in V | uv \in E\}$. The close neighbourhood of a vertex v is $N[v] = N(v) \cup$ $\{v\}$. The order |V(G)| of G is denoted by p. The degree of v is d(v) = |N(v)|. The maximum degree of a graph G is denoted by $\Delta(G)$ and the minimum degree is denoted by $\delta(G)$. A vertex cover in a graph G is a set of vertices that covers all the edges of G. The vertex covering number $\alpha_0(G)$ is the minimum cardinality of a vertex cover in G. A set of vertices in a graph G is called independent set if no two vertices in the set are adjacent. The vertex independence number $\beta_0(G)$ is the maximum cardinality of an independent set of vertices. For a vertex v of a graph G, the eccentricity e(v) is the distance between v and a vertex farthest from v. The maximum eccentricity is its diameter, diam(G). A set D of vertices in a graph G is called a dominating set of G if every vertex in V - D is adjacent to some vertex in D. The domination number of G, denoted by $\gamma(G)$ is the minimum cardinality of a dominating set. The domination in graphs with many variations is now well studied in graph theory. A thorough study of domination appears in [2]. Let S(G) be the subdivision graph of G. The lict graph of n[S(G)] of S(G) is a graph whose vertex set is the union of the set of edges and set of cutvertices of S(G) in which two vertices adjacent if and only if the corresponding members are adjacent or incident. A subset D^d of V[n(S(G))] is double dominating set of n[S(G)] if for every vertex $v \in$ V[n(S(G))], $|N[v] \cap D^d| \ge 2$, that is v is in D^d and has at least one neighbour in D^d or v is in $V[n(S(G))] - D^d$ and has at least two neighbours in D^d . The lict subdivision double dominating number $\gamma_{ddnb}(G)$ is a minimum cardinality of the lict subdivision double dominating set of G and is denoted by $\gamma_{ddns}(G)$. The graph valued function related to domination parameters have been studied in [4,5,6,7,9,10,11,12,13,14,15]. Further in [8], the subdivision of G with graphvalued function and related with domination parameters has been established. In this paper, we establish some sharp bounds for $\gamma_{dans}(G)$. Also some upper and lower bounds on $\gamma_{ddns}(G)$ in terms of the vertices, edges and other different parameters of G and not in terms of the element of n[S(G)]. Further, its relation with other different dominating parameters is also obtained.

II. RESULTS

We need the following theorems to prove our results.

1) Theorem A[1]: For any path
$$P_n$$
, the vertex covering number is $\alpha_0(P_n) = \begin{cases} \frac{n}{2} & \text{if n is even} \\ \frac{n-1}{2} & \text{if n is odd} \end{cases}$



International Journal for Research in Applied Science & Engineering Technology (IJRASET)

ISSN: 2321-9653; IC Value: 45.98; SJ Impact Factor: 6.887 Volume 6 Issue IV, April 2018- Available at www.ijraset.com

- 2) Theorem B[1]: For any path P_n , the edge covering number is $\alpha_1(P_n) = \begin{cases} \frac{n}{2} & \text{if } n \text{ is even} \\ \frac{n+1}{2} & \text{if } n \text{ is odd} \end{cases}$
- 3) Theorem C[2]: A graph G is eulerian if and only if G is of even degree.
- 4) Theorem D[3]: For any connected graph G, with $p \ge 4$, $G \ne K_p$, $\gamma_{ss} = \alpha_0(G)$.

III. THE LICT SUBDIVISION DOUBLE DOMINATION NUMBER OF A GRAPH.

In this section, we characterise the lict subdivision double domination number by giving a necessary and sufficient condition for it and also establish the conditions for a lict subdivision double dominating set

1) Theorem 3.1: A double dominating set D^d of the n[S(G)], is a lict subdivision double dominating set of G, if and only if the following conditions hold. (i) $n[S(G)] - D^d$ has

at least two vertices.

(ii) For any two vertices $u, v \in$

 $V[n(G)] - D^d$, every uv path contains a vertex of D^d .

We now give a characterization of lict subdivision double dominating set of G which is minimal.

2) Theorem 3.2: A double dominating set D^d of the n[S(G)] is minimal if and only if for every vertex $v \in D^d$ either (i) $|N(v) \cap D^d| \le 2$ or (ii) there \exists a vertex $u, \in D^d$

 $V[n(G)] - D^d$ such that $|N(v) \cap D^d| = 2$ and $u \in N(v)$.

Proof: Let D^d be a minimal lict subdivision double dominating set of G. Suppose that there exists a vertex $v \in D^d$ for which $|N(v) \cap D^d| \ge 2$ and for every vertex $u \in V[n(S(G))] - D^d$, either $|N(v) \cap D^d| \ge 2$ or $u \notin N(v)$. Then consider $D^{d'} = D^d - \{v\}$, since v is adjacent to at least two vertices of $D^{d'}$ it follows that $D^{d'}$ is double dominating set of n[S(G)], which is contradicting to the minimality of D^d . Conversely, assume that D^d is double dominating set of n[S(G)] satisfying conditions (i) and (ii). For that consider the set $D^{d'} = D^d - \{v\}$ for any vertex $v \in D^d$. If condition (i) holds, then $|N(v) \cap D^{d'}| \le 2$, which implies that $D^{d'}$ is not a double dominating set of $n[S(G)] - D^{d'}$ such that $|N(u) \cap D^{d'}| = 2$ and $u \in N(v)$. But in this case that set $D^{d'}$ would not double dominating to u and hence would not be a double dominating set of n[S(G)]. Therefore D^d is a minimal double dominating set of n[S(G)].

IV. LOWER BOUNDS FOR $\gamma_{ddns}(G)$.

We establish lower bounds for $\gamma_{ddns}(G)$ in terms of elements of G.

1) Theorem 4.1: For any connected (p, q) graph G with $p \ge 2$, $\left\lfloor \frac{p}{2} \right\rfloor + 1 \le \gamma_{ddns}(G)$. Equality hold if G is P_2 .

Proof: Let $D^d = \{v_1, v_2, \dots, v_n\}$ be a double dominating set of n[S(G)]. By the definition of $n[S(G)], V[n(S(G))] = E(S(G)) \cup C(S(G))$. Let D^d be the double dominating set of n[S(G)] such that any vertex $v \in V[n(S(G))] - D^d, |N[v] \cap D^d| \ge 2$. Then $\{V[n(S(G))] - D^d\}$ contains at least one vertex which gives $\frac{p}{2} < \left[\frac{p}{2}\right] < \left[\frac{p}{2}\right] + 1 \le \gamma_{ddns}(G)$.

2) Theorem 4.2: For any connected (p, q) graph G with $p \ge 2$, $\gamma_{ddns}(G) \le p + q - 1$.

Proof: Let *G* be a (p,q) graph, then V[S(G)] = p + q. Let $F = \{v_1, v_2, ..., v_k\}$ be the set of vertices in S(G) and $F \subseteq V[S(G)]$ such that $|F| = \alpha_0(S(G))$ similarly for $E = \{e_1, e_2, ..., e_n\}$ be the set of edges in S(G), $E \subseteq E[S(G)]$ such that $|E| = \beta_0(S(G))$ and by the definition of n[S(G)], $V[n(S(G))] = E(S(G)) \cup C(S(G))$. Where E(S(G)) is the set of edges and C(S(G)) is the set of cutvertices in S(G). We consider a set $F \subseteq E(S(G))$. If *F* corresponding to such vertices $D^d = \{v_1, v_2, ..., v_q\} \subseteq V[n(S(G))]$ such that $D^d = V[n(S(G))] - C(S(G))$ and any vertex $v \in V[n(S(G))] - D^d$ is dominated by at least two vertices of n[S(G)]. Thus D^d is double dominating set of n[S(G)]. Now consider $|D^d| \le \alpha_0(S(G)) + \beta_0(S(G) - 1 = V[S(G)] - 1 = p + q - 1$. Hence $\gamma_{ddns}(G) \le p + q - 1$.

3) Theorem 4.3: For any connected (p,q) graph G with $p \ge 3$, $p + \gamma(G) \le \gamma_{ddns}(G)$. Equality holds if $G \cong P_{p_i}$, $p \ge 3$.

Proof: Let $D = \{v_1, v_2, ..., v_k\}$ be a minimal dominating set of G such that $|D| = \gamma(G)$. Since in n(S(G)), $V[n(S(G))] = E(S(G)) \cup C(S(G))$. Further let $E = \{e_1, e_2, ..., e_q\}$ be the set of all egdes which are incident to the vertices of D becomes E(S(G)) = 2E, in n(S(G)). Since $V[n(S(G))] = E(S(G)) \cup C(S(G))$. Let $D^d = \{v_1, v_2, ..., v_n\} \subseteq E(S(G)) \subseteq V[n(S(G))]$ be the double dominating set of n(S(G)) such that $|N[v] \cap D^d| \ge 2$. Then D^d form a minimal double dominating set in n[S(G)]. Clearly



International Journal for Research in Applied Science & Engineering Technology (IJRASET)

ISSN: 2321-9653; IC Value: 45.98; SJ Impact Factor: 6.887 Volume 6 Issue IV, April 2018- Available at www.ijraset.com

 $|E(S(G))| - |V(G)| \ge |D|$. Thus it follows that $\gamma_{ddns}(G) - p \ge \gamma(G)$. Hence $p + \gamma(G) \le \gamma_{ddns}(G)$. For equality if $G \cong P_3$. Then in this case |D| = 1, further $S(P_3) = P_5$, $\gamma_{ddns}(G) = 4$. Hence $\gamma_{ddns}(G) = 4 = 1 + 3\gamma(G) + p$.

4) Theorem 4.4: For any connected (p, q) graph G with $p \ge 3$, then $diam(G) + 2 \le \gamma_{ddns}(G)$.

Proof: Let $I = \{e_1, e_2, ..., e_n\} \subseteq E(G)$ be the set of edges which constitutes the longest path between any two distinct vertices of G such that |I| = diam(G). Let $D^d = \{v_1, v_2, ..., v_n\}$ be the set of vertices in V[n(S(G))] such that for any vertex $v \in V[n(S(G))] - D^d$ is adjacent to at least two vertices of D^d and $|N[v] \cap D^d| \ge 2$. It follows that $|D^d| \ge 2$ and the diametral path includes at least two vertices. Thus $diam \le \gamma_{ddns}(G) - 2$ it gives $diam(G) + 2 \le \gamma_{ddns}(G)$.

5) Theorem 4.5: For any connected (p, q) graph G with $p \ge 2$, $\frac{(2q-p(p-3))}{2} \le \gamma_{ddns}(G)$.

Proof: Let $D^d = \{v_1, v_2, ..., v_n\} \subseteq V[n(S(G))]$ be the set of vertices and every vertex $v \in V[n(S(G))] - D^d$ is adjacent to at least two vertices of D^d , thus D^d itself is a double dominating set of n[S(G)]. Since $V[n(S(G))] = E(S(G)) \cup C(S(G))$. Then there exists a vertex $v \in D^d$ which is not adjacent to any vertex of V[n(S(G))]. This implies that $q \leq \frac{p(P-1)}{2} - (p - \gamma_{ddns}(G))$. Which is $2q \leq p^2 - p - 2p + 2\gamma_{ddns}(G)$. Hence the result.

6) Theorem 4.6: For any connected (p, q) graph $G, \left[\frac{2p}{\Delta(G)+2}\right] \leq \gamma_{ddns}(G)$.

Proof: Let D^d be a minimal double dominating set of n[s(G)]. Since in n[S(G)], $V[n(S(G))] = E(S(G)) \cup C(S(G))$. Let k denote the number of edges between D^d and $V[n(S(G))] - D^d$. Since for any connected graph G, there exists at least one vertex $v \in V(G)$ such that deg $(v) = \Delta(G)$, $k \leq \Delta(G) \cdot \gamma_{ddns}(G)$. Also since each vertex v in $V[n(S(G))] - D^d$ is adjacent to at least two vertices in D^d , $k \geq 2(p - \gamma_{ddns}(G))$. From these two inequalities, $2p - 2p\gamma_{ddns}(G) \leq \Delta(G) \cdot \gamma_{ddns}(G)$. It follows that $\frac{2p}{\Delta(G)+2} \leq \left|\frac{2p}{\Delta(G)+2}\right| \leq \gamma_{ddns}(G)$.

7) Theorem 4.7: For any non-trivial tree T of order p with l leaves and s support vertices, $(p - l - s + 4) \le \gamma_{ddns}(T)$.

Proof: To prove that if *T* is a tree of order $p \ge 2$ with *l* leaves and *s* support vertices then $\gamma_{ddns}(T) \ge (p - l - s + 4)$. We use mathematical induction on p. Let p = 2, then $T = P_2$ and diam(T) = 1 it implies that $\gamma_{ddns}(T) = p - l - s + 4 = 2$. Let p = 3 then *T* is a star and diam(G) = 2, it implies that $\gamma_{ddns}(T) = (p - l - s + 4) = 4$. Let p = 4 then *T* may be P_4 and diam(G) = 3 it implies that that $\gamma_{ddns}(T) > (p - l - s + 4)$. For p = n, $\gamma_{ddns}(T) > (p - l - s + n)$. Further it is also true for p = n + 1. Hence we obtain the desired result.

8) Theorem 4.8: For any connected (p, q) graph G with $p \ge 2, \gamma(G) + 1 \le \gamma_{ddns}(G)$. Equality hold if G is K_2 .

Proof: Let $D = \{v_1, v_2, ..., v_k\}$ be a minimal dominating set of G such that $|D| = \gamma(G)$. Since in n(S(G)), $V[n(S(G))] = E(S(G)) \cup C(S(G))$, where E(S(G)) is the set of edges and C(S(G)) is the set of cutvertices in S(G). We consider the following cases. Case 1: Suppose G is a tree. Then clearly for any tree T, $D^d = \{v_1, v_2, ..., v_n\}$ set of vertices in n[S(G)] such that $D^d = V[n(S(G))] - C(S(G))$ is dominated by at least two vertices of n[S(G)]. Since the number of edges of S(G) are more than that of G, which gives $V(G) \subset V[n(S(G))]$. It follows that $|D^d| > |D| + 1$ which gives $\gamma_{ddns}(G) > \gamma(G) + 1$. Case 2: Suppose G is not a tree. Then there exists at least one cycle. Let D be a minimal dominating set of G, then $|D| = \gamma(G)$. Suppose there exists a cycle of length l, then in S(G), the cycle length be 2l. Let $E = \{e_1, e_2, ..., e_k\} \subseteq E[S(G)]$, such that $|E| = \gamma'(G)$. In n[S(G)], $V[n(S(G)] = E(S(G)) \cup C(S(G))$. Suppose $F \subset C[S(G)]$. Then $\{E\} \cup \{F\} \subseteq V[n(S(G))]$ such that $\forall v_i \in V[n(S(G))] - \{E\} \cup \{F\}$ is adjacent to at least two vertices of V[S(G)]. Hence $|E| \cup |F|$ is a minimal double dominating set n[S(G)] such that $|E| \cup |F| = \gamma_{ddns}(G)$ -set. Clearly $|D| \le |E| \cup |F|$ which gives $\gamma(G) + 1 \le \gamma_{ddns}(G)$. For equality, let $G = K_2$, we have $\gamma(G) = 1$. Further $S(K_2) = P_3$ and $n[S(K_2)] = C_3$ and $\gamma_{ddns}(G) = 2$, $\gamma(G) = 1$. Hence $\gamma_{ddns}(G) = 2 = \gamma(G) + 1$. 9) Theorem 4.9: For any connected (p, q) graph G, $\begin{bmatrix} (p(9-p)-6 \\ 6 \end{bmatrix} \le \gamma_{ddns}(G)$.

Proof: Let D^d be a γ_{ddns} -set of G. Since in n[S(G)], $V[n(S(G))] = E(S(G)) \cup C(S(G))$. Let t_1 denote the number of edges in n[S(G)] incident to the vertices of $V[n(S(G))] - D^d$ only. Also t_2 denotes the number of edges in n[S(G)] incident to the vertices of D^d only. Then $\frac{[p(p-1)]}{2} \ge t_1 + t_2 \ge 4|V(G) - D^d| - 2 + |D^d - 1|$ it implies that $\frac{[p(p-1)]}{2} \ge 4p - 4|D^d| - 2 + |D^d| - 1 = 4p - 3|D^d| - 3$. It gives that $3|D^d| \ge \frac{p(9-p)-6}{2}$. Hence $\left[\frac{(p(9-p)-6)}{6}\right] \le \gamma_{ddns}(G)$. 10) Theorem 4.10: For any connected (p,q) graph G with $p \ge 2$, $2(p-q) \le \gamma_{ddns}(G)$.



International Journal for Research in Applied Science & Engineering Technology (IJRASET) ISSN: 2321-9653; IC Value: 45.98; SJ Impact Factor: 6.887 Volume 6 Issue IV, April 2018- Available at www.ijraset.com

Proof: Let $D^d = \{v_1, v_2, ..., v_n\}$ be the minimal set of vertices which covers all the vertices in V[n(S(G))]. Suppose for every vertex of $v \in V[n(S(G))]$ is adjacent to at least two vertices of D^d , clearly D^d forms a double dominating set of n(S(G)). Let any vertex $v \in D^d$ which is not adjacent to any vertex of $V[n(S(G))] - D^d$. Then $2q \ge |D^d| + 2|V(G) - D^d|$ it gives that $2q \ge |D^d| + 2p - 2|D^d|$. Therefore $|D^d| \ge 2p - 2q$ this implies that $2(p - q) \le \gamma_{ddns}(G)$.

V. SPECIFIC VALUES OF $\gamma_{ddns}(G)$.

We found some constraints for which $\gamma_{ddns}(G)$ follows the equality relation with other domination parameters of G.

1) Theorem 5.1: For any path P_p with $p \ge 3$, $\gamma_{ddns}(P_p) = \begin{cases} 4\alpha_0(P_p) - 2, & \text{if } p \text{ is even} \\ 4\alpha_0(P_p), & \text{if } p \text{ is odd} \end{cases}$.

Proof: Let P_p be the path with $p \ge 3$ vertices. Consider $V = \{v_1, v_2, ..., v_n\}$ be the vertices and $E = \{(v_i, v_{i+1})\}$, i = 1, 2, 3, ... be the edge set of path P_p . Since in $n[S(G)], V[n(S(G))] = E(S(G)) \cup C(S(G))$. We consider the following cases

- a) Case I: If p is even. Then by Theorem A[1], we have $\alpha_0(P_p) = \frac{p}{2}$ it implies that $p = 2\alpha_0(P_p)$. Since $\gamma_{ddns}(P_p) = 2q = 2(p 1)$, we have $\gamma_{ddns}(P_p) = 2p 2 = 2[2\alpha_0(P_p)] 2 = 4\alpha_0(P_p) 2$.
- b) Case 2: If p is an odd. Then by Theorem A[1], we have $\alpha_0(P_p) = \frac{p-1}{2}$ it implies that $p 1 = 2\alpha_0(P_p)$. Since $\gamma_{ddns}(P_p) = 2q = 2(p-1) = 4\alpha_0(P_p)$.
- 2) Theorem 5.2: For any path P_p with $p \ge 3$, $\gamma_{ddns}(P_p) = \begin{cases} 4\alpha_1(P_p) 2, & \text{if } p \text{ is even} \\ 4\alpha_1(P_p) 4, & \text{if } p \text{ is odd} \end{cases}$.

Proof: Let P_p be the path with $p \ge 3$ vertices. Consider $V = \{v_1, v_2, ..., v_n\}$ be the vertices and $E = \{(v_i, v_{i+1})\}$, i = 1, 2, 3, ... be the edge set of path P_p . Since in $n[S(G)], V[n(S(G))] = E(S(G)) \cup C(S(G))$. We consider the following cases.

- a) Case 1: If p is even. Then by Theorem B[1], we have $\alpha_1(P_p) = \frac{p}{2}$ it implies that $p = 2\alpha_1(P_p)$. Since $\gamma_{ddns}(P_p) = 2q = 2(p 1)$, we have $\gamma_{ddns}(P_p) = 2p 2 = 2[2\alpha_1(P_p)] 2 = 4\alpha_0(P_p) 2$.
- b) Case 2: If p is an odd. Then by Theorem B[3], we have $\alpha_1(P_p) = \frac{p+1}{2}$ it implies that $p 1 = 2\alpha_1(P_p)$. Since $\gamma_{ddns}(P_p) = 2q = 2(p-1) = 2p 2 = 4\alpha_1(P_p) 4$.

VI. UPPER BOUNDS FOR $\gamma_{ddns}(G)$.

We establish upper bounds for $\gamma_{ddns}(G)$ in terms of elements of G.

1) Theorem 6.1: For any connected (p, q) graph G with $p \ge 2$, $\gamma_{ddns}(G) \le 2q$.

Proof: Since $V[n(S(G))] = E(S(G)) \cup C(S(G))$, V(S(G)) = p + q. Let D^d be double dominating set of n[S(G)]. Then by definition of lict subdivision double domination $|D^d| \ge 2$. Further by definition of n[S(G)], $2q - \gamma_{ddns}(G) \ge 0$. Clearly it follows that $\gamma_{ddns}(G) \le 2q$.

2) Theorem 6.2: For any connected (p,q) graph G with $p \ge 2$, $\gamma_{ddns}(G) \le p + q - \delta(G)$.

Proof: Let $D^d = \{v_1, v_2, ..., v_n\} \subseteq V[n(S(G))]$ be the set of vertices and every vertex $v \in V[n(S(G))] - D^d$ such that $|N[v] \cap D^d| \ge 2$. Thus it is clear that $|D^d| \ge 2$. Since for any graph *G* there exists at least one vertex $v \in V(G)$ such that $\deg(v) = \delta(G)$. By definition of $n[S(G)], V[n(S(G))] = E(S(G)) \cup C(S(G))$. Then there exists a vertex $v \in G$ such that $\deg(v) = \delta(G)$. Thus $\delta(G) \le p - D^d + q$, which implies that $D^d \le p + q - \delta(G)$. Hence the result.

3) Theorem 6.3: For any connected (p,q) graph G with $n[S(G))] \neq K_p$, $\gamma_{ddns}(G) + \gamma_{ssns}(G) \leq 2q + C$. Where C is the number of cut vertices in S(G).

Proof: Suppose *G* has $p \leq 3$ then γ_{ddns} –set does not exist. Now we consider any graph with $p \leq 4$, such that $n[S(G)] \neq K_p$. Since $\gamma_{ddns}(G) \leq 2q$ and from Theorem D[3] $\gamma_{ssns}(G) = \alpha_0[n(S(G))]$. Further $\gamma_{ddns}(G) + \gamma_{ssns}(G) \leq 2q + \alpha_0[n(S(G))] \leq V[n(S(G))] \leq E(S(G)) \cup C(S(G)) \leq 2q + C$. Hence the result.

4) *Theorem 6.4:* For any non-trivial tree *T*, the lict subdivision of a tree is non-eulerian.

Proof: Let *T* be a non-trivial tree and n[S(G)] always contain a point of odd degree. Hence by Theorem C[2], the result follows. NORDHAUS-GADDUM TYPE RESULTS

5) Theorem 6.5: For any connected (p, q) graph G with $p \ge 3$ vertices,



International Journal for Research in Applied Science & Engineering Technology (IJRASET) ISSN: 2321-9653; IC Value: 45.98; SJ Impact Factor: 6.887

Volume 6 Issue IV, April 2018- Available at www.ijraset.com

 $(\mathbf{i})\gamma_{ddns}(G)+\gamma_{ddns}(\bar{G})\leq 4q.$

(ii) $\gamma_{ddns}(G)$, $\gamma_{ddns}(\overline{G}) \leq 4q^2$.

REFERENCES

- [1] Gray Chartrand and Ping Zhan, Introduction to Graph Theory, Tata McGraw-Hill Pub.Co.Ltd. (2006).
- [2] F. Harary, Graph Theory, Addison-Wesley, Reading, Mass, 1972.
- [3] Kulli V.R., Theory Of Domination In Graphs, Vishwa International Publication, Gulbarga, India, 2010.
- [4] T.W. Haynes, S. T. Hedetniemi and P.J. Slater, Fundamentals of domination in graphs. Marcel, Dekker Publ. 1998 New York.
- [5] M.H. Muddebihal and Megha Khandelwal, Strong split lict domination of a graph, International Journal of Engineering and Scientific Research, Vol.2(7), pp 187-198, 2014.
- [6] M.H. Muddebihal and Megha Khandelwal, Strong split block cut vertex domination of a graph, , International Journal of Mathematics and Computer Applications Research, vol.5(5), pp-73-80, 2015.
- [7] M.H. Muddebihal and Megha Khandelwal, On strong split middle domination of a graph, International Journal of recent and innovation trends in computing and communication trends in computing and communication vol.4(6), pp-538-541,2016.
- [8] M.H. Muddebihal and Sumangaladevi, Roman lict domination in Graphs, Ultra Scientist, vol.24 (3) A, pp 449-458, 2012.
- [9] M.H. Muddebihal and Sumangaladevi, Connected roman domination in graph, International Journal of research in engineering and technology, vol.2 (10), pp 333-340, 2013.
- [10] M.H. Muddebihal and Sumangaladevi, Roman block domination in graphs, International Journal of research in science and advanced technologies, vol.6 (2), pp 267-275,2013
- [11] M.H. Muddebihal and Sumangaladevi, Roman subdivision domination in graphs, International Journal of engineering, science and research technology, vol.3(3), pp 1441-1447, 2014.
- [12] M.H. Muddebihal and Suhas P. Gade, Line double domination in graphs, International Journal of Mathematics and Computer Application Research vol.5(3), pp 75-80, 2015.
- [13] M.H. Muddebihal and Suhas P. Gade, Lict double domination in graphs, Global Journal of Pure and Applied Mathematics vol.13(7), pp 3113-3120, 2017.
- [14] M.H. Muddebihal and Suhas P. Gade, Semitotal block double domination in graphs, International Journal of Mathematics trends and Technology, vol.52 (7), 2017.
- [15] M.H. Muddebihal and Suhas P. Gade, Block double domination in graphs, International Journal of Mathematical Archive, vol.9(1), 2018, pp 1-5.











45.98



IMPACT FACTOR: 7.129







INTERNATIONAL JOURNAL FOR RESEARCH

IN APPLIED SCIENCE & ENGINEERING TECHNOLOGY

Call : 08813907089 🕓 (24*7 Support on Whatsapp)